

Alliance: A Complete CAD System for VLSI Design

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Abstract

*The **Alliance** package is a complete set of CAD tools for the specification, design and validation of digital **VLSI** circuits. Beside the tools, **Alliance** includes also a set of cell libraries, from standard cells for automatic place and route tools, to custom block generators to be used in high performance circuits. This package is used in more than 250 universities worldwide.*

*Each **Alliance** tool can operate as a standalone program as well as a part of the complete design framework. After introducing briefly the design methodology, we outline the fonctionnality of the tools. Experimental results conclude the presentation.*

***Alliance** runs on any **Unix** system and has been also ported to Windows using **Cygwin** environment. It is freely available on `ftp`, and includes binaries, leaf cells libraries, on-line documentation, and tutorials.*

1 Introduction

The **Alliance** package is the result of a ten years effort spent at the **LIP6** Laboratory (formerly **MA SI**) of the Pierre et Marie Curie University (UPMC), in Paris. During these years, our major goal was to provide our undergraduate and graduate students with a complete CAD framework, designed to assist them in digital **VLSI CMOS** course. The *Architecture* team at **LIP6** focuses its activity on two key issues: computer architectures using high complexity ASICs, and innovative CAD tools for **VLSI** design. Strong interaction exists between the people working on computer architectures and the one working on CAD tools. The main CAD action aims at fulfilling both the needs of experienced designers by providing practical answers to state-of-the-art problems (logic synthesis, procedural generation, layout verification, test and interoperability), and novice designers, by providing a simple and consistent set of tools. Our **VLSI** design flow is therefore based on both advanced CAD tools that are not available within commercial CAD systems, such as functional abstraction or static timing analysis, and standard design/validation tools.

Alliance VLSI CAD System is free software. Binaries, source code and cells libraries are freely available under the GNU General Public Licence (**GPL**). You are welcome to use the software package even for commercial designs without any fee. You are just required to mention : "Designed with Alliance © LIP6/Université Pierre et Marie Curie". For any questions please mail to : alliance-users@asim.lip6.fr.

1.1 Process independence

To be useful, a CAD system must provide a way to the silicon, therefore **Alliance** provides a large set of cell libraries also available at the layout level. The target technologies of **Alliance** is **CMOS**. The layout libraries rely on a symbolic layout approach that provides process independence in order to allow the designers to easily port their designs from one silicon supplier to another. The main point in this approach is that the pitch matching constraints in both x and y direction are kept through technological retargeting. The translation, fully automated, relies on a technological file suited to a given process.

These files can be generated directly from the process design rules. Also technological files for several processes are available through the CMP and EuroPractice services, provided you signed a NDA for the process.

1.2 Software portability

The **Alliance** package has been designed so to run on an heterogeneous network of workstations. The only requirements are a **C** compiler and a **Unix** system. For the graphical applications, the XWindow library is used. Several hardware platforms, from Intel 386 based computers to Sparc-Stations and DEC Stations, are supported.

1.3 Modularity

According to the interoperability constraints, each **Alliance** tool can operate as a standalone program as well as a part of the complete **Alliance** design framework. Each **Alliance** tool therefore supports several standard **VLSI** description formats : **SPICE**, **EDIF**, **VHDL**, **CIF**, **GDS2**. In that respect, the tools outputs are fully usable under the **Compass** and **Cadence Opus** environnement, provided these tools have the necessary configuration files. The **Alliance** tools support a zero-default top-down design methodology with not only construction tools — layout editor, automatic place & route — but also validation tools, from design rule checker to functional abstraction and formal proof.

1.4 Compactness

Unlike commercially available CAD systems, the **Alliance** CAD Framework suits the limited ressources of low-cost workstations. For small educational projects — 5000 gates —, a **Unix** system with 8 to 20 Mbytes of memory, appropriate disk storage (30 Mbytes per user), and graphic capabilities (X-Window) is sufficient.

1.5 Easiness

All tools and the proposed design flow are simple to teach and to learn. In any situation, easiness and simplicity have been preferred to sophisticated approaches.

To each tool correspond a unique behavior and utility. Each step of the design methodology corresponds to the use of one or a few tools, for which the usage is well identified.

From a pratical point of view, both on-line documentation (**Unix man**) and paper are available with each tool of the **Alliance** package.

2 Alliance design flow

We refer to the term "design flow" as a sequenced set of operations performed when realizing a **VLSI** circuit. In the design flow, we rely on a strict definition of all the objects and design functions found in the process of designing a **VLSI** chip. The design flow is based on the Mead-Conway model and is characterized by its top-down aspect. Below we introduce the major steps of the basic design methodology. It emphasizes the top-down aspect of the design flow, and points out that our methodology is broken up into 5 distinct parts, the latter being not available yet within **Alliance**:

- capture and simulation of the behavioral view,
- capture and validation of the structural view,
- physical implementation of the design,
- layout verification,
- test and coverage evaluation.

The design flow also includes miscellaneous tools like layout editor for the design of the cell libraries, and a PostScript plotter for documentation.

2.1 Capture and simulation of the behavioral view

Like we just saw, the capture of the behavioral view is the very first step of our design flow. Within **Alliance**, any **VLSI** design begins with a timing independent description of the circuit with a subset of **VHDL** behavior primitives. This subset of **VHDL**, called `vbe`, is fairly restricted: it is the data-flow subset of this language. It is not very easy to modelize an architecture using this subset, but it has the great advantage of allowing simulation, logic synthesis and bit level formal proof on the same files.

A **VHDL** analyzer called `vasy` can be used to automatically convert most common **VHDL** descriptions (using for example IEEE 1164 packages) to the restricted Alliance subsets (`vbe` and `vst`).

Patterns, **VHDL** simulation stimuli, are described in a specific formalism that can be captured using a dedicated language `genpat`. Once a **VHDL** behavioral description written and a set of test vectors have been determined, a functional simulation is ran. The behavioral **VHDL** simulator is called `asimut`. It validates the input behavior, according to the input/output vectors.

2.2 Capture and validation of the structural view

The structural view can be captured once the data flow description is validated. The actual capture of the netlist relies either on specific description languages, `genlib` for standard cells or `dpngen` for data-path, or on direct synthesis from the data flow using the `boom` tool for optimization and the `boog` tool to map on a cell library. `Genlib` and `dpngen` are netlist-oriented libraries of C functions. In the design methodology, it is essential for the students to get acquainted with the C language basics. The advantage of such an approach is that designers do not have to learn several language with specific syntax and semantics.

Usually, the main behavior is partitionned in several sub-behaviors. Some are described recursively using the `genlib` language and the other ones can be directly synthesized from a **VHDL** description of the corresponding sub-behaviors. The `boog` tool takes an **RTL** description and generates a netlist of standard cell gates. An other subset of **VHDL** allows to capture finite state machines.

This subset, called `fsm`, can be translated into a **RTL** description using the tool `syf`, and then the resulting description optimized using `boom` and finally synthesized as a netlist using once more `boog`.

Since `asimut` can operate on both **RTL** and structural views, the structural description is checked against the behavioral description by using the same set of patterns that has been used for behavioral validation.

2.3 Physical design

Once the circuit netlist has been captured and validated, each leaf of the hierarchy has to be physically implemented. A netlist issued from `boog` is usually placed and routed using the over cell router `nero`. If the netlist has been captured using `genlib` and if it has a high degree of regularity, it can be placed manually for optimisation using other `genlib` functions.

The different parts can be placed and assembled together using `ocp` and routed using overcell router called `nero`, and this generates what we call a *core*. The circuit core is now ready to be connected to external pads. The core-to-pads router, `ring`, aims at doing this operation automatically, provided the user has given an appropriate netlist and some indications on pad placement.

The last stage of the physical implementation is the translation of the symbolic layout to a foundry compliant layout using the `s2r` tool. After that, the tape containing the circuit can be processed by the silicon supplier.

2.4 Verification

In our **VLSI** class, we intend to show that **VLSI** verification is at least as important as **VLSI** physical design. For that reason, we have introduced in the design flow powerful tools to perform behavior, netlist and layout verifications.

The correctness of the design rules is checked using the design rule checker `druc`.

An extracted netlist can be obtained from the resulting layout. `Cougar`, the layout extractor operates on both hierarchical and flattened layout and can output both flattened netlists (transistor netlist) and hierarchical netlists. The transistor netlist can be the input of a spice simulator.

When extracted hierarchically, the resulting netlist can be compared with the original netlist by using the `lvx` tool. `lvx`, that stands for Logical Versus Extracted, is a netlist comparator that matches every design object found in both netlists.

The critical path of the circuit is evaluated using a commercial static timing analyzer, as **Alliance** doesn't provided one.

2.5 Test and coverage evaluation

For now, the fault coverage provided by the functional patterns is evaluated using a commercial fault simulator, as **Alliance** doesn't provide one yet.

3 Tools and layout libraries of the Alliance package

Every **Alliance** tool has been designed to simply interface with each other, in order to support the proposed design flow. Nevertheless, each tool can also be used independently, thanks to the multiple standard formats used for input and output files.

One of the most important characteristics of the **Alliance** system is that it provides a common internal data structure to represent the three basic views of a chip:

- the behavioral view,
- the structural view,
- the physical view.

Figure 1 details how all the **Alliance** tools are linked together around the basic behavioral, structural and physical data structures.

The process independence goal is achieved with a thin fixed-grid symbolic layout approach. All the library of the system use this approach successfully. Layouts have been targetted to ES2 $2\mu\text{m}$, $1.5\mu\text{m}$, $1.2\mu\text{m}$, $1.0\mu\text{m}$ and $0.7\mu\text{m}$ technologies, the AMS $1.2\mu\text{m}$ technology and SGS-Thomson $0.5\mu\text{m}$, $0.35\mu\text{m}$ technologies. Chips have been fabricated successfully through the **CMP** services on these technologies.

3.1 Tools

- `vasy` is a **VHDL** analyzer and convertor. The supported **VHDL** subset is closed to commercial synthesizer tools such as Synopsys. It converts automatically **VHDL** descriptions to the restricted **VHDL** subsets of Alliance tools.
- `asimut` is a **VHDL** logic simulator. The supported **VHDL** subset allows both structural and behavioral data-flow description (without timing information). Complex systems and microprocessors, including **INTEL** 8086 and **MIPS** R3000 have been successfully simulated with `asimut`. `Asimut` is based on an event-driven algorithm and powerful representation of boolean functions using binary decision diagrams.
- `genpat` is a language interpreter dedicated to efficient descriptions of simulation stimuli. It generates an **ASCII** file that can act as an input of `asimut`. A `genpat` file format to **MSA** translator allows the generation of appropriate simulation patterns for the Tektronix LV500 tester. This allows to perform functional tests when the circuits comes back from the foundry.
- `loon` is a gate level netlist optimizer. If the output of the logic synthesis takes into account the internal delays of the cells during the mapping phase, it doesn't take into account the fan-out problems. `Netoptim` work is to ensure that the drive capabilities of all cells are correct, and to try to minimize the delays on the critical pathes in inserting buffers where appropriate.
- `genlib` is a procedural language for netlist capture and placement description (there is no schematic editor in the **Alliance** system). `Genlib` provides a consistent set of **C** primitives, giving the designers the ability to describe **VLSI** circuit netlists in terms of terminals, signals and instances, or circuit topologies in terms of placement of abutment boxes. `Genlib` is mainly used to build parameterized netlist and layout generators.
- `dpngen` is a language that has moreorless the same functionalities as `genlib`, but it is dedicated to datapath description. Its primary difference with `genlib` is that it allows to manipulate vectors of cells, like 32 two inputs `nand` gates or a 32 bits adder. It contains many primitives that greatly simplify the description of operative parts, in an optimized manner. `dpngen` has been recently merged with `genlib`.
- `boom` is a logic optimizer and logic synthesis tool. The input file is a behavioral description of the circuit using the same **VHDL** subset as the logic simulator. The boolean equations described in **VHDL** are optimized so to minimize the number of boolean operators. The output is a new, optimized, data flow description.



Figure 1: How the tools are linked on the data structures.

- `boog` is a logic synthesis tool. The output is a netlist of gates. `boog` can map a data-flow description on any standard-cell library, as long as a **VHDL** data-flow description is provided with each cell.
- `syf` is a finite state machine synthesizer. More precisely, `syf` assigns values to the symbolic states used for the automaton description, and aims at minimizing the resulting logic for both state transition and output generation. The input is a `fsm` description, using a dedicated subset of **VHDL** that includes process description. The output is a behavioral description of the circuit using the same **VHDL** subset as the logic simulator. The output of `syf` is to be synthesized into a netlist of gates using `boog`.
- `ocp` is a placer for standard-cells. The placement system is based on simulated annealing.
- `nero` is an over cell router. The input is a netlist of gates and a placement file. The output is an hierarchical chip core layout without external pads. A specialized router is used for core to pad routing.
- `Ring` is a specific router dedicated to the final routing of chip core and input/output pads. `Ring` takes into account the various problems of pad placement optimization, power and ground distribution. A set of symbolic pads is included in the package.
- `S2r` is the ultimate tool used in our design flow to perform process mapping. `S2r` stands for "symbolic to real", and translates the hierarchical symbolic layout description into physical layout required by a given silicon supplier. The translation process involves complex operations such as denotching, oversizing, gap-filling and layer adaptation. Output formats are either **CIF** or **GDSII**. `S2r` requires a parameter file for each technology aimed at. This file is shared with `druc`, `cougar`, `graal` and `dreal`. From an implementation point of view, these tools use a bin-based data-structure that has very good performances.
- `druc` is a design rule checker. The input file is a - possibly hierarchical - symbolic layout. It checks that a layout is correct regarding the set of symbolic design rules. This correctness must be ensured in order for `s2r` to produce a layout compatible with the target silicon foundry.
- `Cougar` is a layout extractor. The input is a - possibly hierarchical - layout. The layout can be either symbolic or real. The output is an extracted netlist with parasitic capacitances and optionally resistors. The resulting netlist can either be hierarchical or flattened (up to transistor level netlist).
- `Lvx` is a logical versus extracted net-compare tool. The result of a run indicates if the two netlist match together, or if there are different. Note that `lvx` doesn't work at the transistor level.
- `proof` performs a formal comparison between two data flow **VHDL** descriptions that share the same register set. `Proof` supports the same subset of **VHDL** as `asimut`, `boom`, `boog`.
- `graal` is an hierarchical symbolic layout editor. It requires a X-Window graphical environment and the Motif libraries. `Graal` is used for cell layout design or hierarchical block construction. It provides an on-line **DRC** and automatic display of equipotential nets. Editing a cell under `graal` is shown figure 2.
- `dreal` is a real layout editor (rectangles in micron) . It requires a X-Window graphical environment and the Motif libraries.
- there are many other tools not described here.



Figure 2: Editing some custom layout using graal.

3.2 Cell libraries

The **Alliance** package provide a wide range of libraries, either static, ie. fixed cells, or dynamic, as the block is produced by running a parameterized generator. These libraries are compatible with any two metals/one polysilicon technology.

Each object in the library has, for static ones, or produces, for dynamics ones, three views at least :

- the symbolic layout, that describes the cell topology.
- the netlist, in terms of transistor interconnections.
- the behavior, specified in **VHDL** data flow form.

3.2.1 Standard cell library

The `sxlib` library contains boolean functions, buffers, mux, latches, flip-flops, ... (around 100 cells). All cells have the same height and N times the width, where N is the number of pitches. That is the only physical information given in the cell list below. Power supplies are in horizontal ALU1 and have the same width. Connectors are inside the cells, placed on a 5x5 grid. Half layout design rules are a warranty for any layer on any face, except for the power supply and NWELL. Cells can be abutted in all directions whenever the supply is well connected and connectors are always placed on the 5x5 grid. They are supposed to be used with a usual standard cells place and route tools, such as **Alliance's** `ocp` and `nero`, **Compass** or **Cadence**. These cells are to be used primary for glue logic, since optimized operators can be obtained using dedicated generators, as stated paragraph 3.2.2. The `logic` tool can map a behavioral VHDL onto this library.

3.2.2 Datapath libraries: (This part of the document is not up to date !)

There are two kinds of datapath libraries:

- `dp_sxlib` is a cell library dedicated to high density data-paths. It must be used in conjunction with the data-path tool `dpgen`. The cells in `dp_sxlib` have the same fonctionnalities as the ones in `sxlib`, but have a topology that is usable only within a datapath. `Boog` can also map a behavior onto the `dp_sxlib` library.
- `fplib` is a set of above 30 regular functions that are useful in the design of a datapath. These functions range from a n inputs nand gate to a n times m register file.

Here the cells share the power and ground lines in metal2. A powerful dedicated over the cell router can route custom blocks and logic glue in the same structure. Among the `fplib` functionalities, four optimized blocks generators should be presented in more details, as they reflect the quality of this library. All the generators are build with a tiler using a dedicated leaf cell library. Their output is a symbolic layout, a **VHDL** behavior, a set of pattern for test purpose, a netlist, an icon, and a datasheet indicating size and timing estimation for a given technology. The structural parameters varies according to their functionalities.

- optimized generators for datapath operators:

`rsa`, a fast adder generator, with propagation time in $\log nb$ and size in $nb \log nb$, where nb is the number of bits. Its has 2 or 3 input buses, and if needed a carry input. It may be used as a subtractor or adder/subtractor.

| Params | Meaning | Range |
|------------------|--------------------|---------------|
| <code>nb</code> | number of bits | 3 to 128 |
| <code>cin</code> | carry in | true or false |
| <code>csa</code> | three inputs adder | true or false |
| <code>ovr</code> | overflow flag | true or false |

`rfg`, a static register file generator. It has one write address, and one or two read address. It may be operated as a set of level-sensitive latches or edge triggered flip-flops.

| Params | Meaning | Range |
|------------------------|--------------------------|--------------------|
| <code>nb</code> | number of bits | 2 to 64 |
| <code>nw</code> | number of words | 2 to 256 |
| <code>bus</code> | number of read bus | 1 or 2 |
| <code>op</code> | mode of operation | latch or flip-flop |
| <code>low power</code> | reduce power consumption | true or false |

`bsg`, a barrel shifter generator. Possible operations are :

- logical right shift
- arithmetical right shift
- logical left shift
- arithmetical left shift
- right rotation
- left rotation

| Params | Meaning | Range |
|-----------------|----------------|---------|
| <code>nb</code> | number of bits | 3 to 64 |

`amg`, an integer modified booth algorithm array multiplier. the x and y inputs are independent, and pipeline stages can be inserted in the circuit.

| Params | Meaning | Range |
|-----------------|---|---|
| <code>nx</code> | number of bits of the x operand | 8 to 64 |
| <code>ny</code> | number of bits of the y operand | 8 to 64 |
| <code>ps</code> | number of pipeline stages to be inserted in the circuit | 0 to $\min(\frac{nx}{2}, \frac{ny}{2})-1$ |

3.2.3 Custom libraries

Two full-custom parameterized generators are also available. They produce stand-alone blocks, that are to be routed only at the floorplan level with other blocks, using either `bbt` or better `xcheops`.

- ROM and RAM generators:

`grog`, a generic ROM generator. The interface is an address bus, a clock and an output enable signal, and a data out bus. The coding format to specify the ROM contents is a limited subset of VHDL.

| Params | Meaning | Range |
|--------|------------------|--|
| nb | number of bits | 1 to 64 |
| nw | number of words | 64, 128, 256, $n \leq 512, 1 \leq n \leq 8$ |
| hz | tri-state output | true or false |

`rage`, a RAM generator. The interface has a read/write address, a write signal indicating if a read or a write is to be performed, and a clock.

| Params | Meaning | Range |
|--------|---|-------------------------|
| nb | number of bits | 2 to 128 |
| nw | number of words | 128 to 4096 |
| aspect | aspect ratio | narrow, medium or large |
| ud | unidirectional, ie share the same bus for data in and out | true or false |

All these generators have been designed using the **Alliance** CAD tools, for both design and verification phases.

3.2.4 Pad library

Alliance provides also a `padlib` library. This library also uses a symbolic layout approach, and therefore a whole chip can be targeted on several technology without even the core to pad routing. A very robust approach has been enforced, as the pads are subject to electrostatic discharge, and also more sensible to latch-up than the other parts of the circuit due to the amount of current that flows through them.

Chips using these pads have been fabricated on ES2 1.0 μ m, AMS 1.2 μ m and SGS-Thomson 0.5 μ m technology and work as expected.

4 Supported exchange formats

The **Alliance** CAD system handles many file formats. They are summarized here. A file can be either read, using a *parser*, or written, using a *driver*.

- Behavioral view:
 - dataflow **VHDL** parser and driver.
- Structural view:
 - **VHDL** parser and driver.
 - **EDIF** parser and driver.
 - **Spice** parser and driver.
 - **Compass** parser and driver.
 - **Alliance** parser and driver.

- **Hilo** driver
- Physical view:
 - **Alliance** parser and driver, for symbolic layout.
 - **Compass** parser and driver, for symbolic layout.
 - **Modgen** parser and driver, for symbolic layout.
 - **CIF** parser and driver, for real layout.
 - **GDSII** parser and driver, for real layout.

Being able to understand and write many file formats is a must. First, in a development environment, as it allows to check the validity of tools on other CAD systems. Second, because some tools are not available or desirable within **Alliance**, but may be useful however: it is possible to feed an other software with a design in that situation.

The experience shows that many of these formats are used daily. For example, the design that we fabricate through the CMP services are transmitted using the **GDSII** format. The final **DRC** on these files are performed using **Cadence** `pdverify`.

An other example: **Alliance** does not have a fault simulator yet. However this kind of tool is very useful to evaluate the fault coverage of a set of vectors and must be introduced in a **VLSI** class. This is hopefully easily done using the **Hilo** output of **Alliance** that feed the `hifault` simulator.

5 Alliance internal organization

The complete **Alliance** CAD system contains about 600 000 lines of C code, and over 600 leaf cells. It compiles and runs on most **Unix** system, and requires the basic X-Window library X11 plus Motif. The distribution tape shows that there are three kinds of files:

- common data structures and manipulation primitives.
- parsers/drivers to read and write external file formats.
- actual tools.

Alliance as been developed in order to simplify cooperative work between the CAD tool designers. The existence of a common data structure framework releaves the developer of many burdens: reading and writing many file format, conceptualizing the VLSI objects, writing classical high level and nevertheless complex functions, ... All the **Alliance** tools share these data structures and their related functions. So each tool communicates with the other ones smoothly, by construction.

6 Use of Alliance inside our laboratory

Alliance is used for both educational and research purposes. We relate our experience below.

| Project | transistors | Functionality |
|------------------|-------------|---|
| Smal | 17 000 | one bit processor for SIMD architectures |
| Data-safe | 35 000 | dynamic data encryption chips |
| TNT | 60 000 | switch-router for T800 transputerss |
| FRISC | 200 000 | floating-point RISC microprocessor |
| StaCS | 875 000 | Very Long Instruction Word processor |
| Rapid2 | 650 000 | SIMD systolic and associative processor |
| Rcube | 350 000 | Message router for parallel machines |

Figure 3: Various chips designed with **Alliance**.

Educational aspects

The **Alliance** System has been extensively used during the past fifteen academic years (1989-2004) as a practical support of two undergraduate courses: one on **CMOS VLSI** design, the other one on **advanced computer architecture**. These initiation courses lasts 13 weeks with a 2 hours lecture and 4 hours spent using the **Alliance** system per week, and involves 60 students and 3 teachers.

The ‘VLSI design’ course is for students that have no previous knowledge on **VLSI** design and mainly come from two distinct channels: "computer science" and "electrical engineering" masters of sciences. During this course, students are required to design and implement an **AMD2901** compatible processor, starting from a commercial data-sheet. The chip, with a complexity of about 2000 transistors, is designed by groups of 2 or 3 students. The main interest in this course is to teach the design methodology. Most of the **Alliance** tools are used during this class.

The ‘architecture’ course focuses on the way processor architecture, from the system point of view and not from an implementation one. Typical **CISC** and **RISC** processors are studied, and part of them modeled using our **VHDL** subset. In that class, only the `asimut` simulator is used.

Alliance is also used in an intensive graduate course, for the design of the 32 bits microprocessor **MIPS RISC** processor – 30000 transistors –. This course lasts two months, and aims only at the implementation : the high level behavioral model of the processor is given to the students. During that period of time, all the **Alliance** tools are used.

Research projects

These projects range from medium complexity ASICs developed in 6 months by a couple of designers **Data-safe**, **TNT**, **Smal**, **Rf264**, etc... to high complexity circuits (**FRISC**, **Multick**, **StaCS**, **Rapid2**, **Rcube**) developed by a team of PhD students.

The three largest circuits described in table 3 use not only standard-cells but also parameterized generators for regular blocks like *RAMs*, data-paths, or floating-point operators. The **FRISC** and **TNT** projects successfully used the **Cadence** and **Compass** place and route tools, and therefore prove the interoperability of the **Alliance** system.

A picture of the **StaCS** processor is shown figure 4.

7 Conclusion

We are very satisfied to use a set of tools of our own for teaching **CMOS VLSI** design for two good reasons. First, we simply can’t afford 50 high end workstations to run commercial CAD systems like **Synopsys**, **Mentor Graphics** or **Cadence**. Second, both the **Synopsys** and **Cadence** system

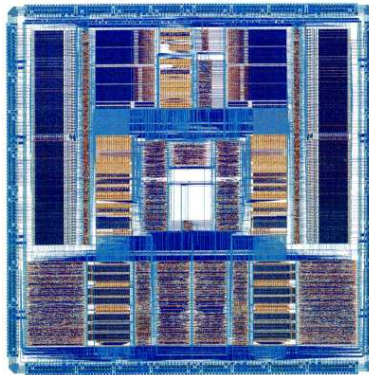


Figure 4: The 875 000 VLIW StaCS processor.

have been used in research project at **LIP6**. They are powerful and sophisticated environments but are much too complex for novice undergraduate students. The great advantage of the **Alliance** CAD system is that we have done our best to stick to the basic yet powerful concepts of **VLSI** design. To each tool correspond a unique fonctionnality, that cannot be changed or worked around by parameter files. At last, we experienced that the technology migration and process independence are key issues. Hence, it is crucial to rely on a portable library at the symbolic layout level.

The **Alliance** package is now in use all over the world, and more than 250 sites have registered today. It is available through anonymous ftp at <ftp://ftp-asim.lip6.fr/pub/alliance/distribution/>, or through a Web browser at <http://www-asim.lip6.fr/pub/alliance/distribution/>.

There is an **Alliance** mailing list, where users can share their views and problems, and our team is always ready to answer questions. The address of this mailing list is alliance-users@asim.lip6.fr.